



#### Tools of the Trade

JONES OF LITTLE I PAGE

5 Shooly Slice to write explain in Cl.

5 are all tools with a python knot short to exploit transforment article

6 Cediting Consociation street.

1 Upher region methors while implements store.

1 Programmed with formed mining tool.

1 clastic large/for agreed was agreed or refutate.

#### Randomness Attacks Mitigation

REINCHIMPOSS ALUBIANS FULURIAL WIFE A MINISTER ARROYS TO A WARRINGTON OF A WIFE A MINISTER ARROYS TO BE WARRING AND A WARRING AN

#### Related Work

#### Stefan Esser

• mt\_srand and not so random numbers.

#### Samy Kamkar

How I met your girlfriend.

Gregor Kopf

Non obvious bugs by example,

#### Summary

- Randomness attacks affect a very large number of PHP applications.
- Exploit mitigations are needed for these attacks,
   Crypto bugs are becoming a trend for exploitation.

#### Thank You!

Questions?

PRNG: Pwning PseudoRandom number Generators (in PHP applications)

George Argyros Aggelos Kiayias

#### PRNG:

### Pwning Random Number Generators (In Php) Applications)

George Argyros\*

Aggelos Kiayias

University Of Athens

\*Census Inc.

# A Small Challenge

```
<?php
   echo mt_rand();
   bet = GET['bet'];
   if (\$bet == mt_rand()) {
      system($_GET['cmd']);
?>
```

# Would you put this code on your server?

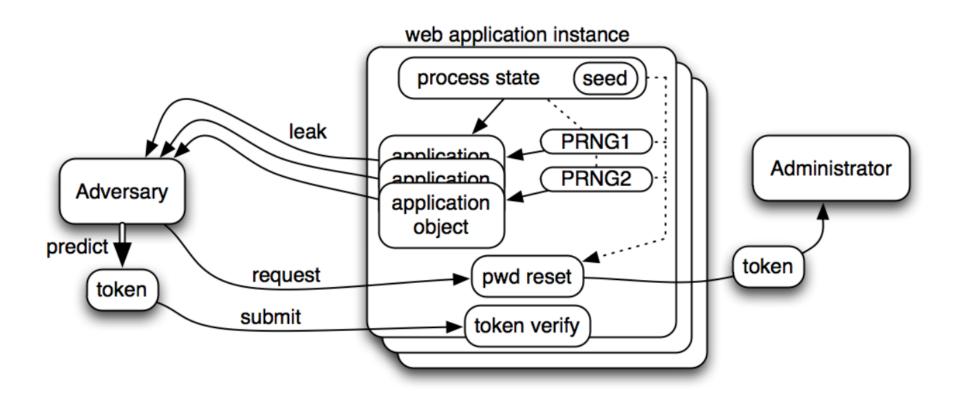
COUC OII YOUI SCIVCI:

Can you exploit this code on a target server?

#### Goal:

Provide practical techniques for the exploitation of randomness vulnerabilities in PHP applications

#### Attack Template



### Entropy Generation in PHP

# PHP Core

#### mt\_rand()/mt\_srand():

- PRNG based on the Mersenne Twister generator.
- mt\_srand() seeds the generator with a 32 bit value
- mt\_rand() produces a 32 bit but PHP truncates the LSB.
- Can return a number in a smaller range.

#### rand()/srand()

- Uses the rand() function from the OS
- seeded with a 32 bit value and produces 31 bit outputs.
- Can return a number in a smaller range.

#### php\_combined\_lcg()/lcg\_value()

- php\_combined\_lcg() is used internall by the PHP system.
- lcg\_value is its public interface.
- Combines two 32 bit LCGs and produces a 64 bit output.
- Seeded only once the first time it is called.

#### uniqid(prefix, extra\_entropy)

- When called without arguments produces a timestamp with seconds and microseconds since unix epoch.
- The first argument adds a user supplied prefix to the timestamp.
- The second argument adds an output of php\_combined\_lcg() as suffix.

### PHP extensions

#### openssl\_random\_pseudo\_bytes()

- Interface to the openssl function with the same name.
- Returns a number of pseudorandom bytes.
- A flag is used to tell if the bytes are crypto strong.
- Requires openssl extension.

#### mcrypt\_create\_iv()

- Gathers entropy from the operating system generators
  - /dev/random, /dev/urandom
- Requires mcrypt extension.

### The Entropy Of Time Measurements



Timestamps make really good PRNG seeds, provided you're willing to kill everyone who owns a clock.

-- Matthew Green

#### Timestamps Facts:

• Epoch (time up to seconds accuracy) is leaked to the client

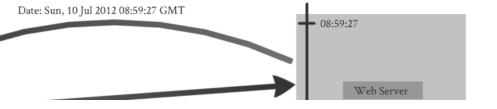
#### Timestamps Facts:

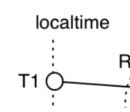
- Epoch (time up to seconds accuracy) is leaked to the client through the HTTP Data Header.
- Microseconds range from 0 to 10<sup>6</sup>
  - Therefore a trivial bruteforce will succeed after 500k requests on average.

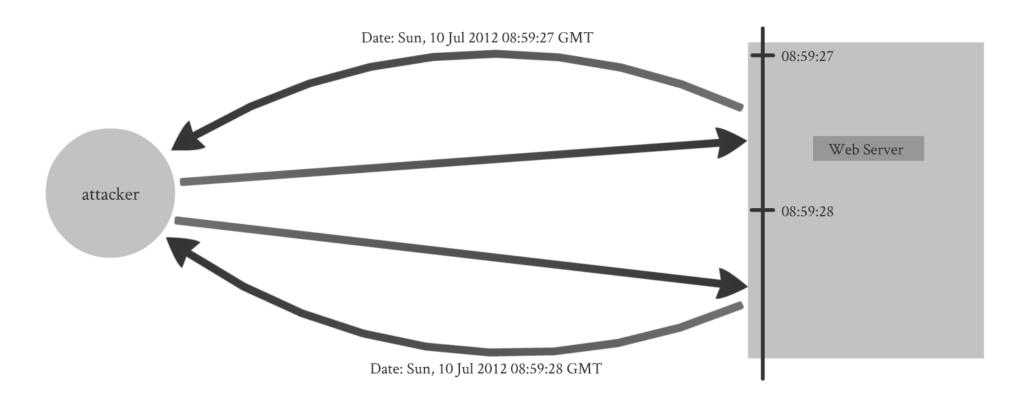
#### Can we do better than 500k?

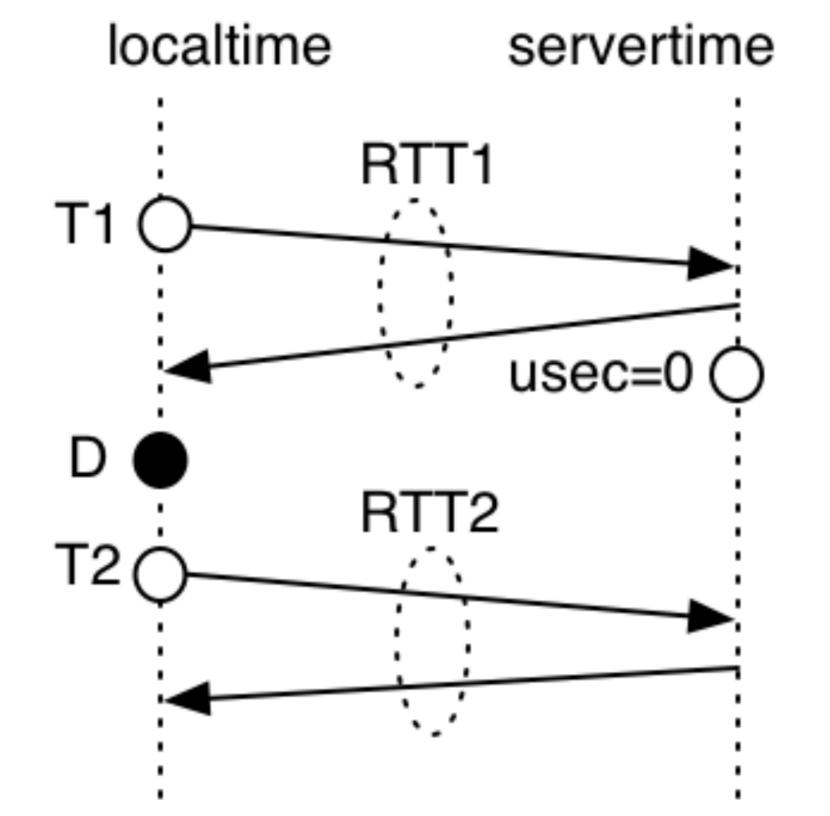
- Adversial Time Synchronization
- Request Twins

#### Adversarial Time Synchronization (ATS)

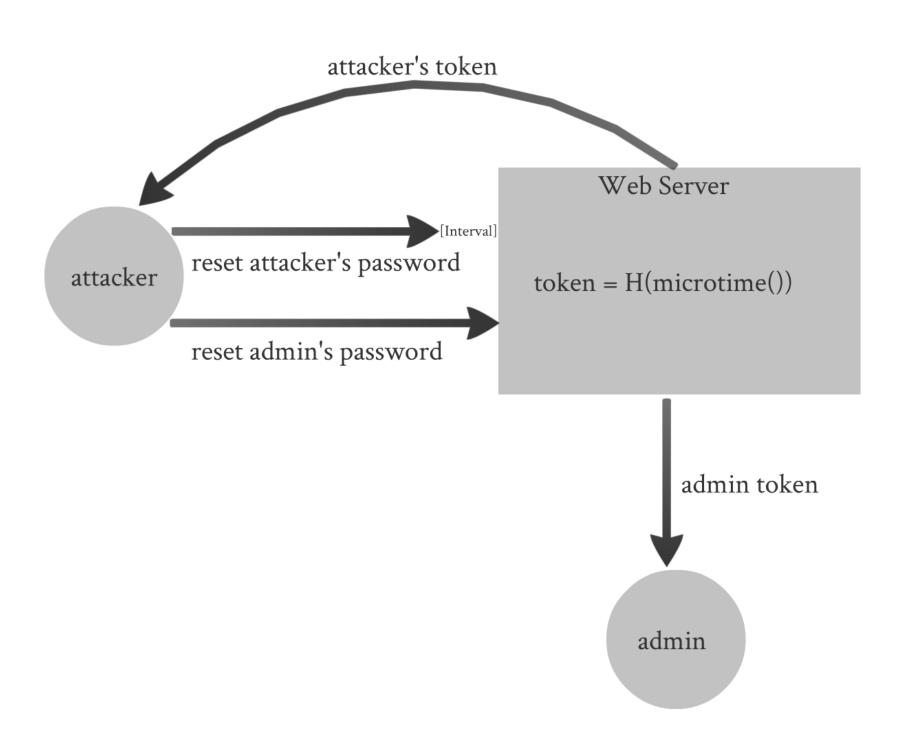








# Request Iwins



# Evaluation

#### Experiment

Predict the output of the following script:

```
<?php
  echo microtime();
?>
```

#### Results

Configuration		ATS			Req. Twins		
CPU(GHz)	RTT(ms)	min	max	avg	min	max	avg
$1 \times 3.2$	1.1	0	4300	410	0	1485	47
$4 \times 2.3$	8.2	5	76693	4135	565	1669	1153
$1 \times 0.3$	9	53	39266	2724	1420	23022	4849
$2 \times 2.6$	135	73	140886	83573	2	1890	299

Time is in microseconds

### Case study



## zencart

The art of e-commerce

- Token generation:
  - seeds mt\_rand with microtime().
  - Produces token using mt\_rand().
- Configuration:
  - 2 cores system
  - RTT = $\sim 10 \text{ ms}$

#### Results:

On average 7k requests to compromise the application

#### bonus:

code execution on the server!

## Attacking a PRNG



# Seed Attacks

In order to attack the seed we need the ability to interact with newly seeded generators. This usually happens when a fresh process is created.

### Process management



### Apache handler (mod\_php):

• PHP runs as an Apache web server module.

#### CGI:

• There is a new PHP process spawned for each request and terminated after the request is served.

#### Fast CGI:

• There are a number of PHP processes serving requests repeatedly. They are usually killed after they served a predefined number of requests.

### Keep-Alive Requests

- When the Connection HTTP header is set to Keep-Alive the web server keep the connection open.
- There is a maximum number of keep-alive requests.
- In mod\_php all requests within the same connection are handled by the same process.
  - Multiple requests to the same PHP process.

### Generating fresh PHP processes

- In mod\_php when the number of occupied processes reaches a certain threshold the server creates new processes to handle subsequent requests.
  - The default threshold in Apache is to have less than 5 idle processes.
- We can exploit this functionality to force the creation of new PHP processes!

### Technique

- Create a large number of connections using the Keep-Alive HTTP header.
- While keeping these connections alive make a new connection to the server.
- The new connection is very likely to be handled by a fresh PHP process.

Hacking your own PHP session identifier

session\_id = MD5(client IP address || time\_of\_day() || php\_combined\_lcg())

- If the total entropy is "small" then we can obtain a preimage by bruteforce.
- This gives us the value of php\_combined\_lcg().

# But what does small means?

- Since we have access to the MD5 sum we can perform the bruteforcing on the CPU rather over the network.
  - 250\$ GPU --> 2^{30} MD5 / sec.
  - 750\$ GPU --> 2^{32} MD5 /sec.
- Entropies up to 40 bits are easily handled...

session\_id = MD5(client IP address || time\_of\_day() || php\_combined\_lcg())

Known since its the attacker's IP address.

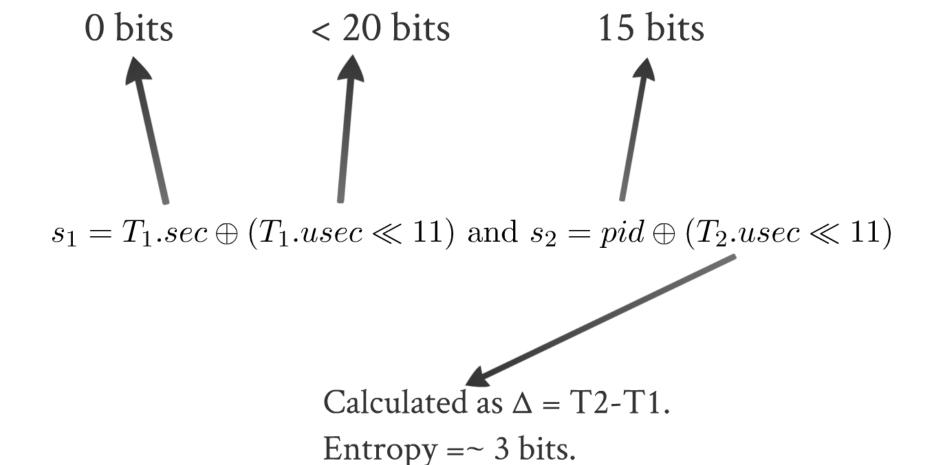
Provides up to 20 bits of entropy which can be reduced using the ATS algorithm.

php\_combined\_lcg has a 64 bits output so bruteforcing the output is not feasible.

Since we can generate fresh processes we can try to predict the first output which is simply one round of the generator with the seed. php\_combined\_lcg has two 32 bit registers s1, s2. Seeding:

 $s_1 = T_1.sec \oplus (T_1.usec \ll 11)$  and  $s_2 = pid \oplus (T_2.usec \ll 11)$ 

- T1, T2: two subsequent timestamps
- pid: PHP process identifier.

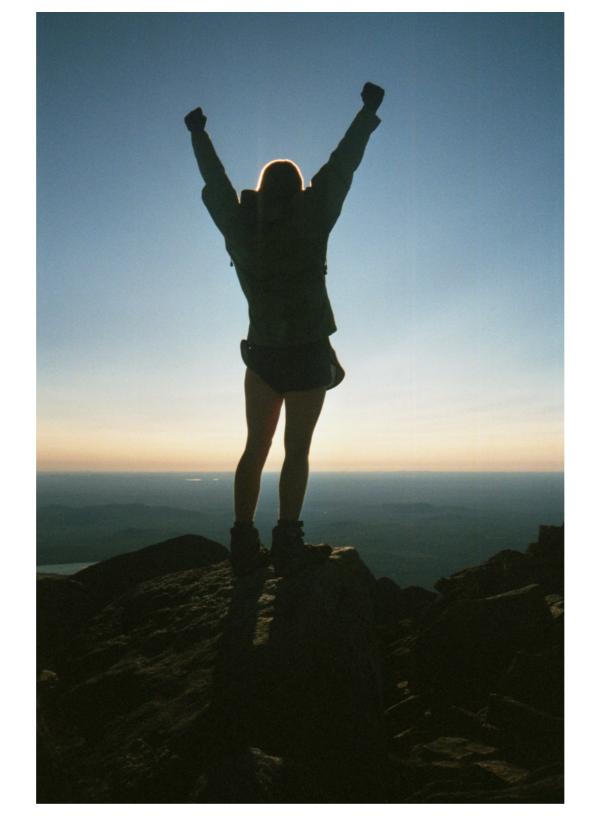


Total entropy of session identifier in a fresh process is about 40 bits.

By bruteforcing the session identifier of a fresh process we can obtain:

- the seed of php\_combined\_lcg
- the process identifier of the PHP process.



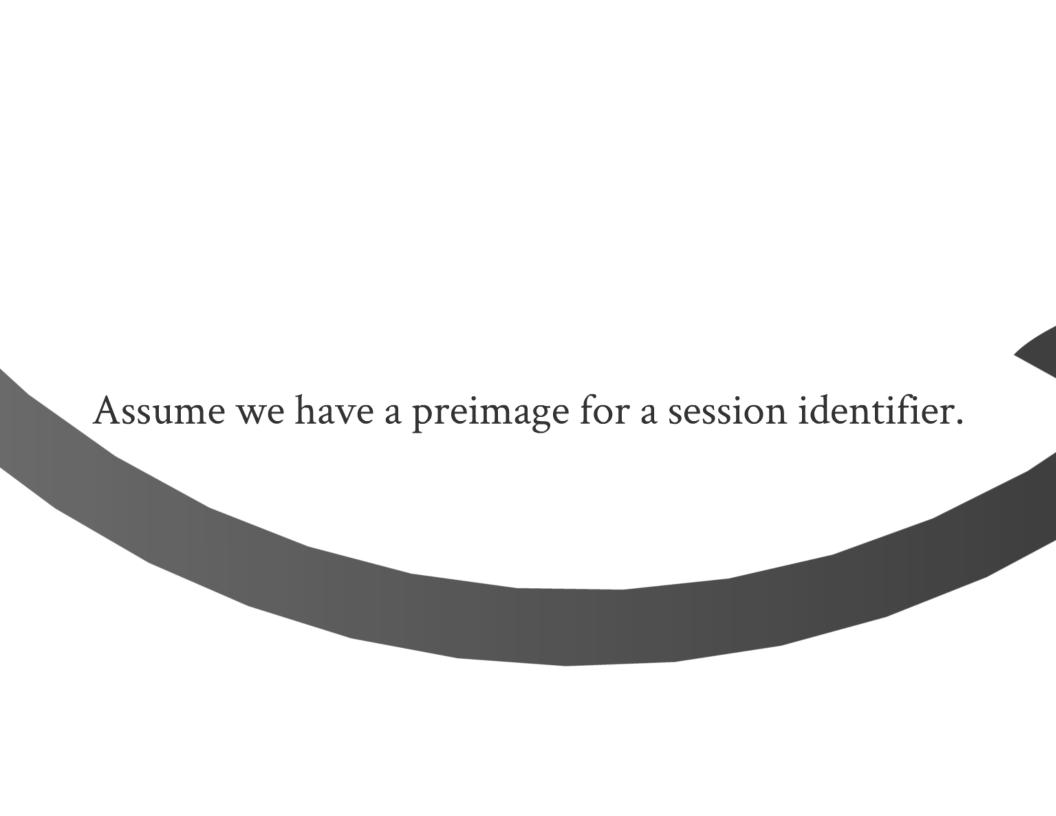


# What about the other PRNGS?

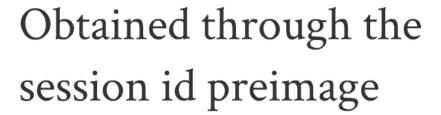
### Seeding in rand()/mt\_rand()

- These generators can be seeded with the respective functions mt\_srand() and srand().
- If the generator is not seeded, then the following 32 bit seed is produced:

$$seed = (time()*pid) ^ (10^{6}*php\_combined\_lcg())$$



Leaked by the Date HTTP Header





 $seed = (time()*pid) ^ (10^{6}*php\_combined\_lcg())$ 

Obtained through the session id preimage

A session identifier preimage completely determines the seed of the mt\_rand() and rand() PRNGs!

The attack does not require any outputs from the targetted PRNGS!

## Case study



\$user->hash = random::hash()

```
\phiuser->\piasm = ramuom...masm()
```

```
static function hash($entropy="") {
  return md5($entropy . uniqid(mt_rand(), true));
}
```

#### Attack:

- Obtain a preimage for a session identifier
  - This will give uniqid's extra entropy and mt\_rand
- Use request twins to bruteforce uniqid's timestamp.



The bruteforce idea of session\_id also applies on the seed of rand() and mt\_rand().

Assume we are connected to a fresh process and obtain some outputs are dependent on rand()/mt\_rand().

Then we can do an offline bruteforcing of all 2^{32} possible seeds to find which one generated the observed outputs.

Contrary to the previous attacks that relied on information obtained online this attack relies only on the small size of the seed.

Therefore the process can be further optimized using an application specific rainbow table.

# Case study



### Password reset algorithm in Joomla

- 2008: mt\_rand() seeded with microtime().
- 2010: mt\_rand() seeded with the crc32 of an unpredictable string along with an installation time generated key produced the same way.
- 2011: Seeding removed, and default PHP seeding was use along with the secret key.



```
$registry->set('secret', JUserHelper::genRandomPassword(16));
public static function genRandomPassword($length = 8)
     $salt = "abcdef[...]QRSTUVWXYZ0123456789";
     $len = strlen($salt);
     $makepass = ";
     for (\$i = 0; \$i < \$length; \$i++)
          $makepass .= $salt[mt_rand(0, $len - 1)];
     return $makepass;
 function getHash($seed)
   return md5(JFactory::getConfig()->get('secret') . $seed);
```

### Notice:

If the configuration script was executed on a fresh PHP process then the entropy of the secret key is 32 bits regardless of its length!

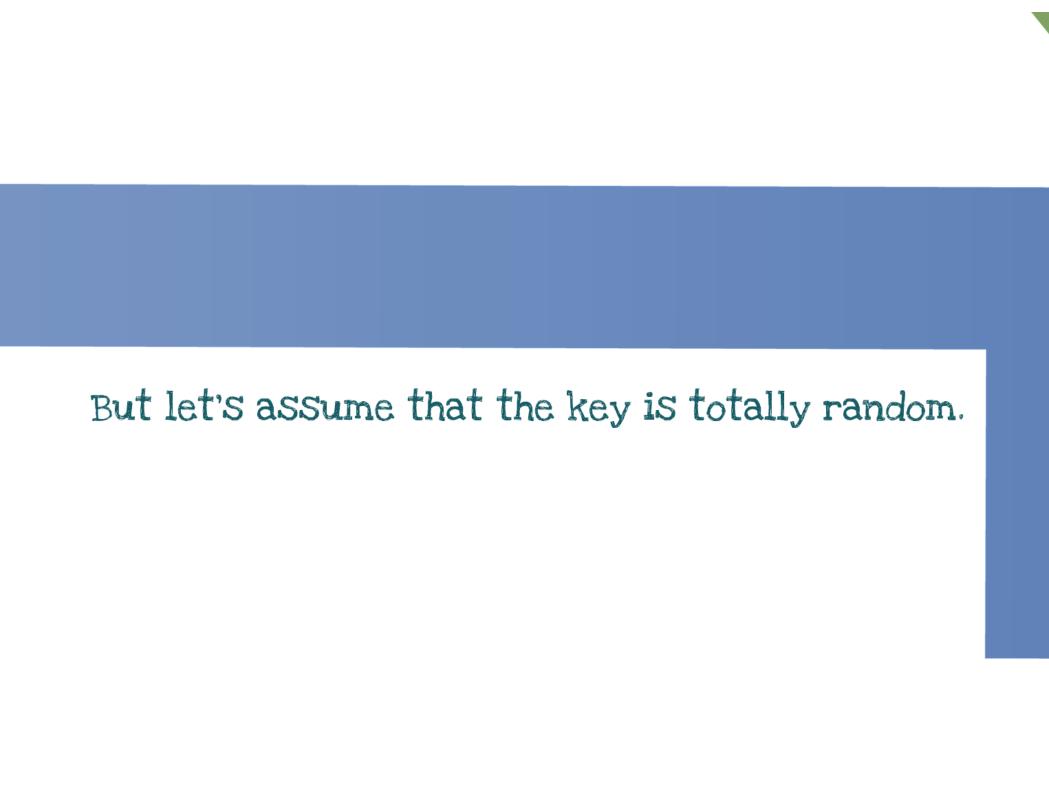
secret key is used for "remember me" cookies.

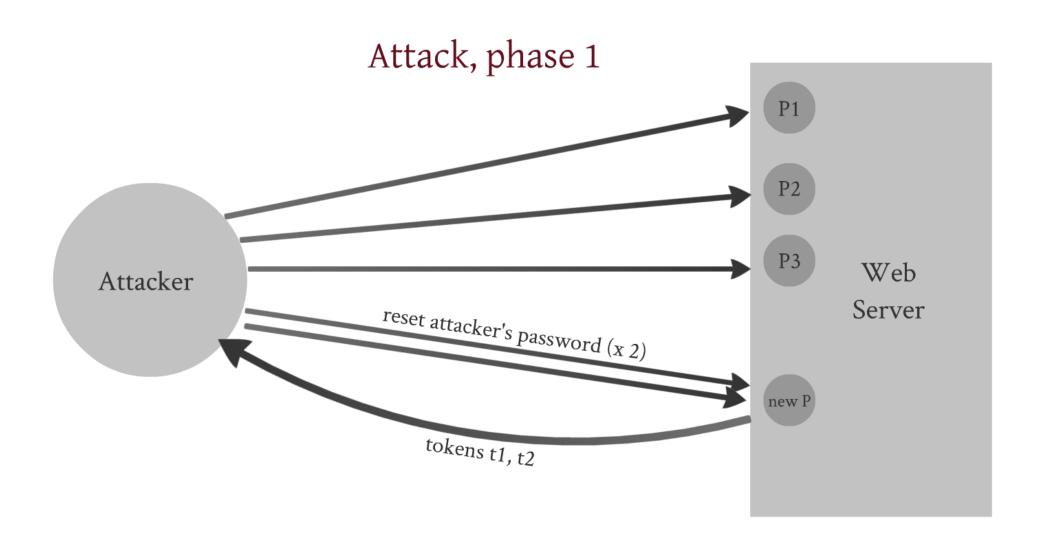
setcookie(self::getHash('JLOGIN\_REMEMBER'), \$rcookie, ...);

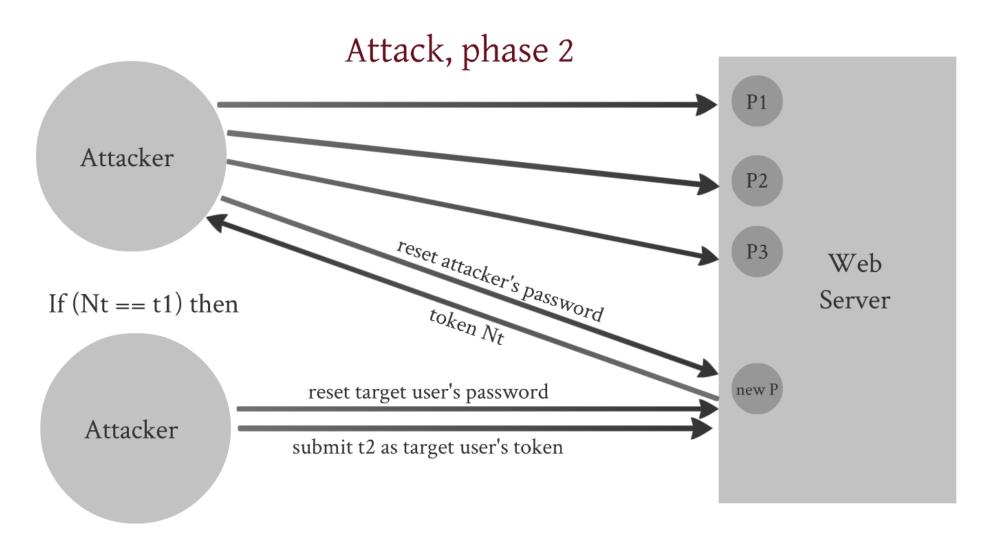
## Vulnerability no.1

- Get a "remember me" cookie and bruteforce all 2^{32} possible secret keys.
- Once the secret key is recovered use one of the previously mentioned attacks to predict future generated tokens.

The attack works only when the secret key is generated from a fresh process.







Expected Number of requests: 2<sup>32</sup>:-(

We can request more than one token in the first phase of the attack.

```
1 token --> 2^{32} requests
2 tokens --> 2^{32}/2 requests
k tokens --> 2^{32}/k requests
```

Total number of requests as a function of the token pairs we will request:

$$F(x) = 2x + 2^{32} / x$$



2 requests for

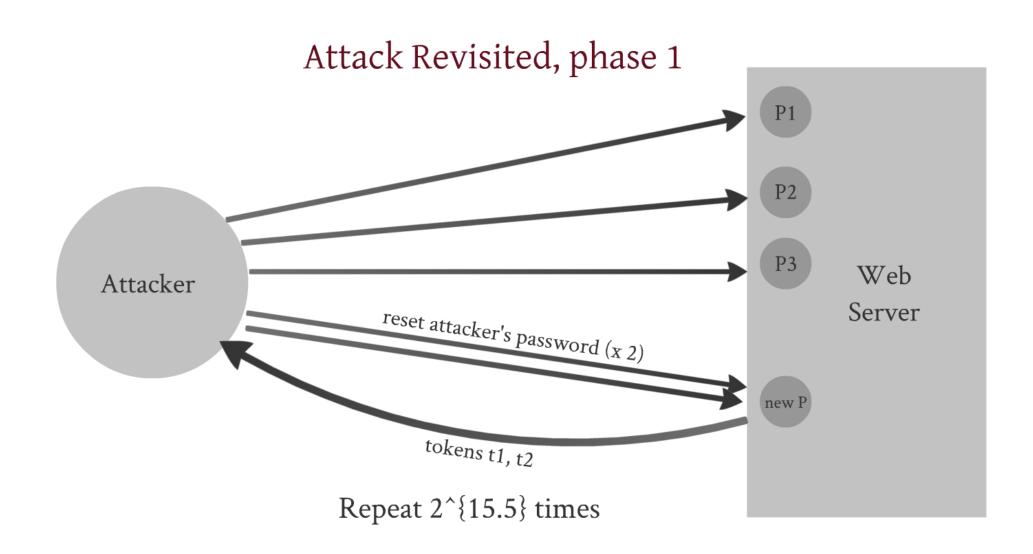
2 tokens

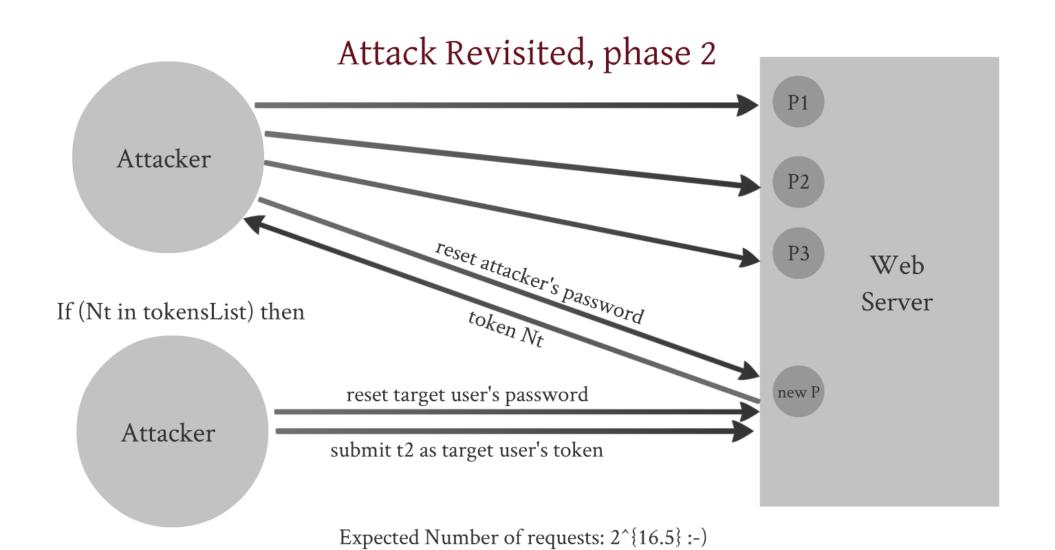


number of requests to hit one of the tokens we have obtained.

Our goal is to minimize function F.

- F minimizes at  $x = 2^{15.5}$  for which we have that F(x) = 185k requests.
- New process creation will incur a 10% overhead in default Apache installations.





This attack shows that a general class of otherwise secure token generation algorithms are vulnerable due to the insecure seed of PHP.

ex. \$token = AES(\$very\_random\_key, mt\_rand().mt\_rand()).



# 수호신

# 

- Hardening extension for PHP.
- Replace rand() with a Mersenne Twister generator with different state than mt\_rand().
- Disables srand() and mt\_srand() functions.
- Seeds the generators once at process startup with a secure seed.
  - Entropy gathered from the operating system.

# Attacking a PRNG



# State Recovery Attacks

# All PRNGs in the PHP core are linear

# Challenges in predictability

## Truncation:

- The output may be truncated to a smaller range.
- This may introduce non linearity to the generator.

# Truncation in PHP

In order to truncate a number n from  $[M] = \{0,...,M-1\}$  to a range[a,b] PHP does the following:

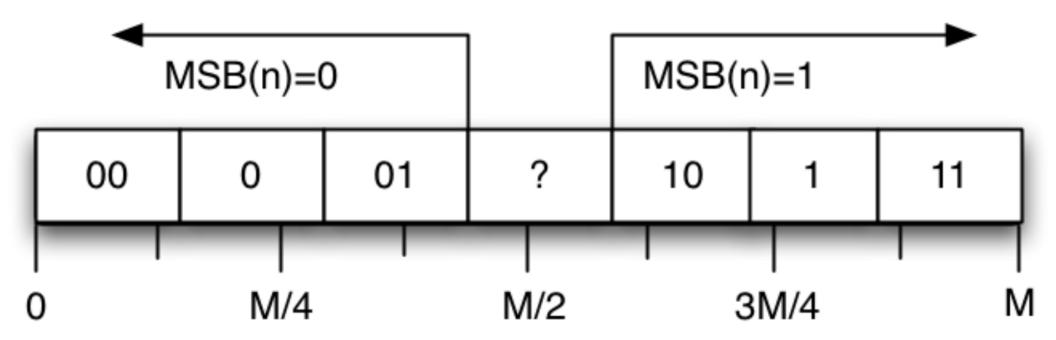
$$l = a + \frac{n \cdot (b - a + 1)}{M}$$

We can view this process as one that puts M values into b-a+1 buckets based on their MSBs.

Given a bucket number we can determine a range for the original number n

$$\left\lfloor \frac{(l-a) \cdot M}{b-a+1} \right\rfloor \le n \le \left\lfloor \frac{(l-a+1) \cdot M}{b-a+1} \right\rfloor - 1$$

Depending on the number of bits common in the upper and lower bound we can determine some of the MSBs of the original number.



# Challenges in predictability

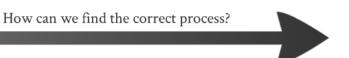
### Process identification:

- We want to get all outputs from the same generator.
- Because of the Keep-Alive limit the server might close the connection before the necessary leaks are collected.

# eaks are collected.

# Process Distinguisher

nit we will be disconnected number of requests.





Due to the Keep-alive limit we will be disconnected from our process after a number of requests.

However we can try to reconnect afterwards.

### Algorithm:

- Connect to a fresh process and obtain a session id prein
- · Obtain DRNC leaks until the correct closes the connecti

How can we find the correct process?

Need access to some process specific state

## Idea:

Use the session id preimage as a process specific state to distinguish between server processes.

### Algorithm:

- Connect to a fresh process and obtain a session id preimage.
- Obtain PRNG leaks until the server closes the connection.
- Start reconnecting to the server and requesting session identifiers.
  - For each session check if it is generated using the next round of php\_combined\_lcg than the one used in the preimage we have.
  - If a match is found then we have connected to our process.

# state recovery for rand()

### rand() implementation depends on the OS:

> On windows a 15 bit LCG is used

$$X_{n+1} = (aX_n + c) \bmod m$$

> On \*nix systems an additive feedback generator is used:

$$r_i = (r_{i-3} + r_{i-31}) \bmod 2^{32}$$

> Truncation introduces non linearity.

#### Hastad-Shamir Framework

Solves the problem of uniquely solving an underdetermined system of linear modular equations when part of the variables is known.

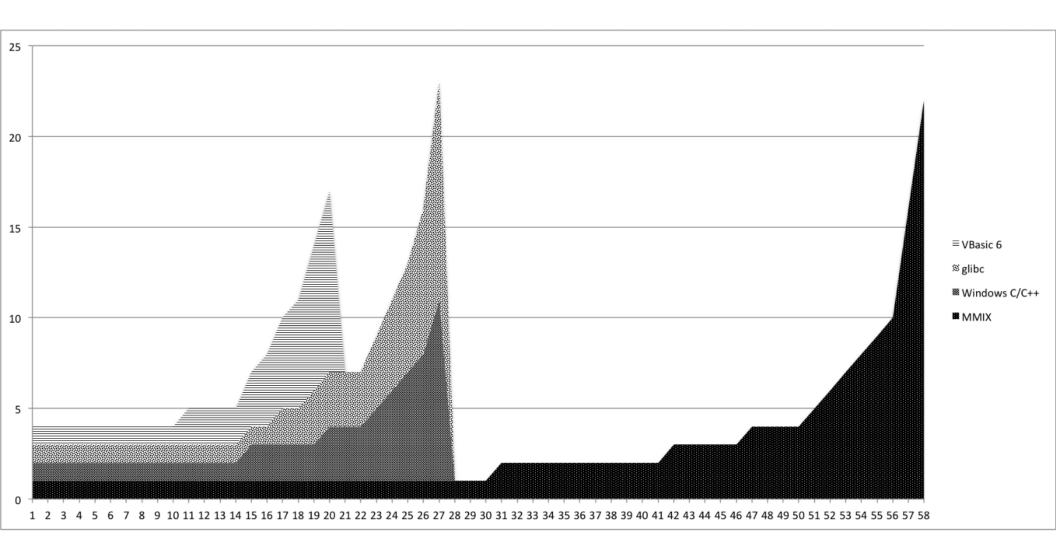
### High level description

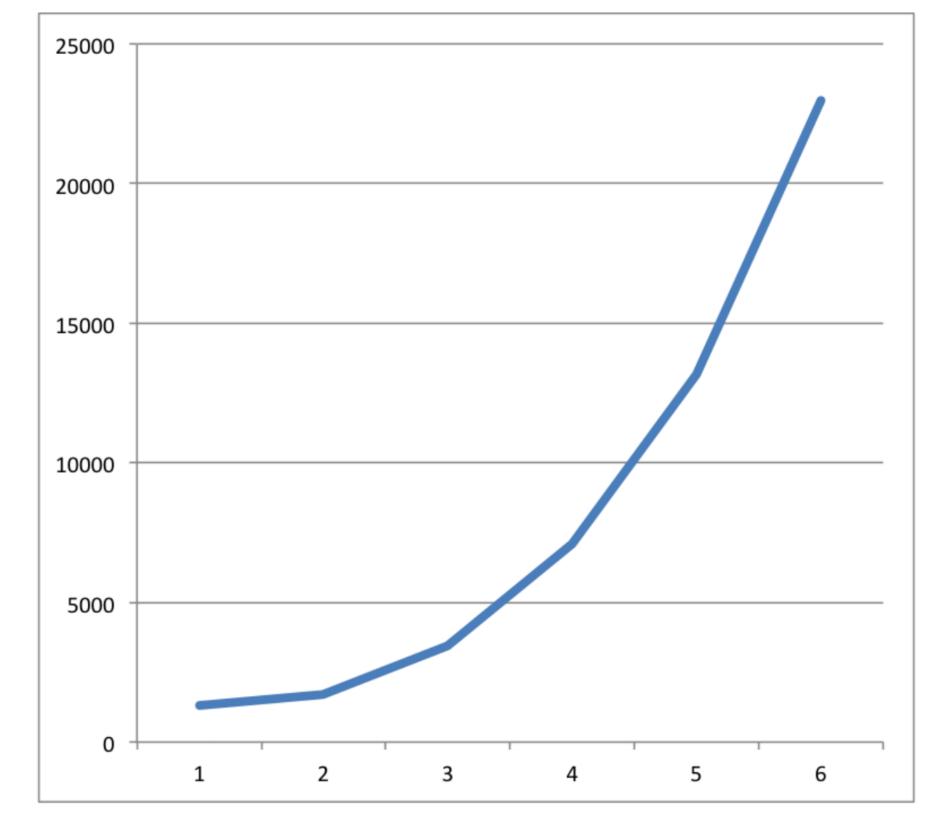
- Define a lattice over the coefficients of the equations
- Find a reduced basis of the lattice using the LLL algorithm.
- Use the fact that the basis vectors are small to uniquely solve the system over the integers.

### Bottleneck point:

- Lattice base reduction of a lattice with dimension equal to the number of leaks needed.
- Public LLL algorithm implementations have complexity  $O(d^5)$ .
  - A O(d^3 logd) variant exists, but without any public implementation.

# Implementation experiments





### Summary:

- LCGs can be efficiently recovered unless we have extremely large truncation levels.
- Similar behavior is observer in the glibc generator however the LLL complexity did not allowed to recover more than 6 bits.

# State recovery for mt\_rand()

#### Notice:

Truncation does not introduce non linearity to the generator.

- Each output bit can be express linear equation to the internal
- Take each output bit obtained create a linear system.
- · If the system has a unique solu

#### Mersenne Twister:

#### Based on a linear recurrence over GF(2):

$$x_{k+n} = x_{k+m} \oplus ((x_k \land 0x80000000) | (x_{k+1} \land 0x7fffffff))A$$

$$xA = \begin{cases} (x \gg 1) & \text{if } x^{31} = 0\\ (x \gg 1) \oplus a & \text{if } x^{31} = 1 \end{cases}$$

Huge state of 19937 bits. (a Mersenne Prime)

To improve randomness properties the output of the recurrence is multiplied by an invertible matrix:

$$z = xT$$

T is called the tempering matrix.

### Notice:

Truncation does not introduce non linearity to the generator.

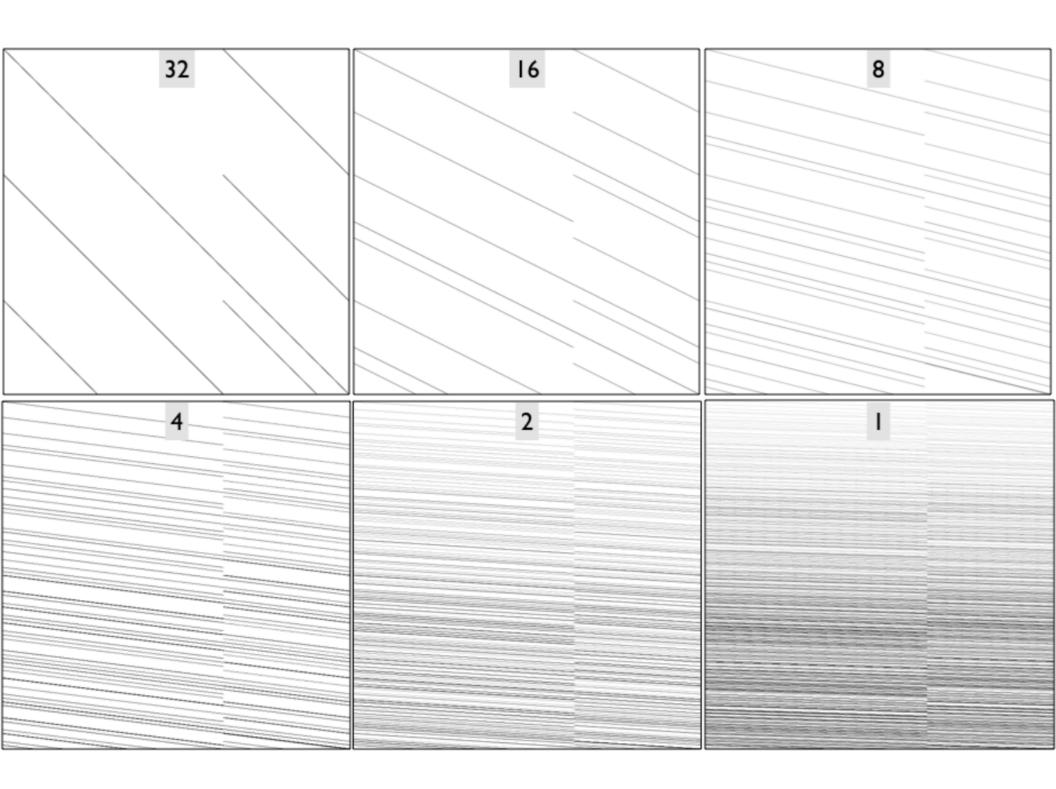
- Each output bit can be expressed as a linear equation to the internal state.
- Take each output bit obtained and create a linear system.
- If the system has a unique solution the solution will give the internal state of the generator.

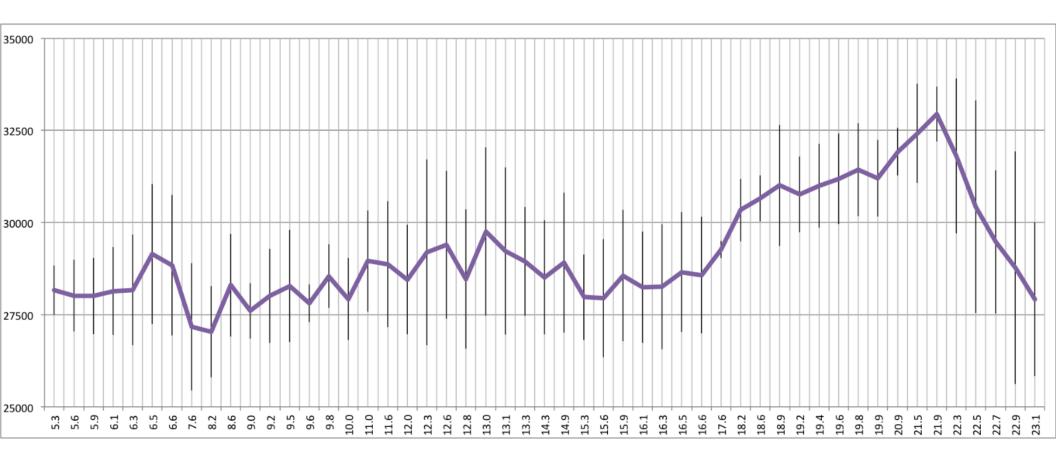
How can we know that the system has a unique solution in advance?

### Employ an online gaussian solver:

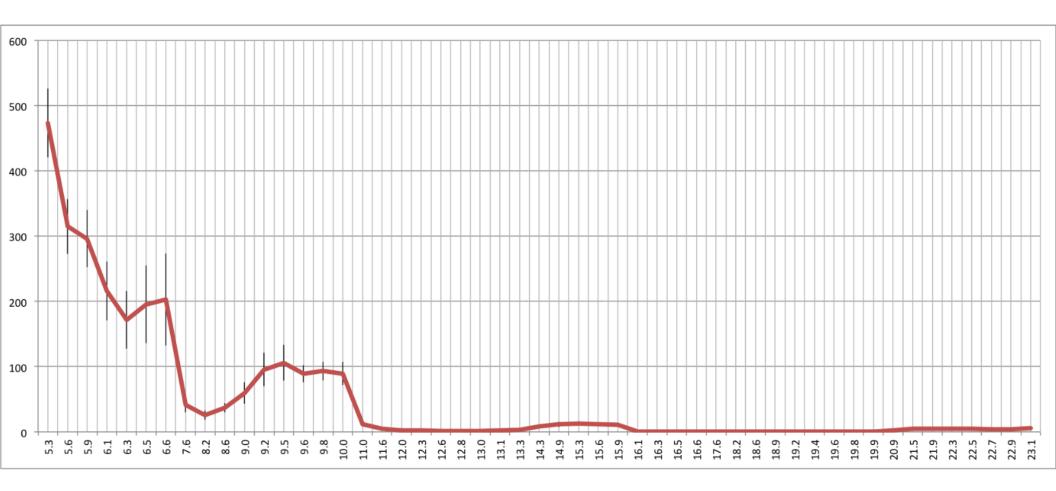
- As equations are obtained from the server add them to the system.
- Stop when the system becomes uniquely solvable.

# Implementation experiments





Equations vs truncation.



Time vs Truncation.

# Case study



```
function phorum_gen_password($charpart=4, $numpart=3)
 vowels = ... //[char array];
 cons = ... //[char array];
 $num_vowels = count($vowels);
 $num_cons = count($cons);
 $password="";
 for(\$i = 0; \$i < \$charpart; \$i++){
   $password := $cons[mt_rand(0, $num_cons - 1)]
              . $vowels[mt_rand(0, $num_vowels - 1)];
 $password = substr($password, 0, $charpart);
 if($numpart){
   $max=(int)str_pad("", $numpart, "9");
   $min=(int)str_pad("1", $numpart, "0");
   $num=(string)mt_rand($min, $max);
 return strtolower($password.$num);
```

At least 4 mt\_rand() outputs skipped in each call.

• The resulting system is very dense.

#### Results:

### On average:

- 11 reconnections of the client.
- 30 minutes to compromise the application.

### Tools of the Trade

- Nobody likes to write exploits in C!
- A set of tools with a python interface in order to exploit randomness attacks:
  - Online Gaussian solver.
  - Lightweight rainbow tables implementation.
  - · Programmable web bruteforcing tool.
- check http://crypto.di.uoa.gr for a release.

## Randomness Attacks Mitigation

- PHP 5.4 added extra entropy to the session identifier.
  - session.entropy\_length enabled by default.
- Suggested to add secure seeding to all PRNGs in the PHP core.
  - PHP security team: "This is an application specific problem".
  - Secure PRGs from extensions are rarely used right now.
- A drop in replacement for any token generator can be found in http://crypto.di.uoa.gr
  - Checks for crypto strong PRGs in the PHP system
  - Otherwise collects entropy from various sources.

### Related Work

### Stefan Esser

mt\_srand and not so random numbers.

### Samy Kamkar

How I met your girlfriend.

### Gregor Kopf

Non obvious bugs by example.

#### Summary

- Randomness attacks affect a very large number of PHP applications.
- Exploit mitigations are needed for these attacks.
- Crypto bugs are becoming a trend for exploitation.

### Thank You!

Questions?

PRNG: Pwning PseudoRandom number Generators (in PHP applications)

George Argyros Aggelos Kiayias